

BOUNDED RATIONALITY AND STATISTICAL PHYSICS

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THE GOLDEN RULE

DO NOT:

*Consider a variable x ,
that optimizes a function*

INSTEAD:

*Consider a distribution over x ,
that optimizes an expectation value*

PROBABILITY COLLECTIVES (PC)

The golden rule gives an underlying language - *Probability Collectives* - for translating between many fields:

- i) Bounded rational game theory**
- ii) Statistical physics (mean field theory)**
- iii) Adaptive control**
- iv) Optimization, constrained or not, over *any* measurable space**
- v) Reinforcement learning**
- vi) Sampling of distributions**

Especially suited to distributed applications

ROADMAP

1) *Review information theory*



2) *Show bounded rationality = statistical physics*



3) *What is distributed control, formally?*



4) *Optimal distributed control policy*



5) *How to find that policy in a distributed way*

REVIEW OF INFORMATION THEORY

- 1) Want a quantification of how “uncertain” you are that you will observe a value i generated from $P(i)$.**
- 2) Require the uncertainty at seeing the IID pair (i, i') to equal the sum of the uncertainties for i and for i'**
- 3) This forces the definition**

$$\text{uncertainty}(i) = -\ln[P(i)]$$

REVIEW OF INFORMATION THEORY - 2

4) So expected uncertainty is the *Shannon entropy*

$$S(P) \equiv -\sum_i P(i) \ln[P(i)]$$

- Concave over P, infinite gradient at simplex border

5) *Information* in P, $I(P)$, is what's left after the uncertainty is removed: $-S(P)$.

6) This allows us to formalize Occam's razor:

Maxent: Given $\{E_P(g_i) = 0\}$, “most plausible” P is the P consistent with $\{E_P(g_i) = 0\}$ having minimal $I(P)$

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***IN THE REAL WORLD, EVERYONE IS
BOUNDED RATIONAL***

- **Real players (human or otherwise) are bounded rational, due to limited computational power if nothing else.**
- **Previous attempts to address this are mostly ad hoc models of (human) players**
 - **Underlying problem of arbitrariness of those models.**

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- **Science: Without information concerning a system, you cannot infer anything concerning it. So ...**

Inference of players' strategies *must* be based on observed/provided information.

COMBINING INFORMATION THEORY AND GAME THEORY

- **Say our information is $P_{(i)}$, the strategies of players other than i , and i 's expected cost.**

Then the minimum information principle says it is “most conservative” to infer that P_i minimizes

$$L_i(P) = \beta E(h_i) - S(P)$$

where S is the Shannon entropy, and β a constant.

- **Alternative: If information is the entropy of i 's mixed strategy, predict that P_i minimizes i 's expected cost:**

Again, P_i minimizes $L_i(P)$

QUANTIFYING BOUNDED RATIONALITY

- At Nash equilibrium, each P_i separately minimizes

$$E(h_i) = \int dz h_i(z) \prod_j P_j(z_j)$$

- Allow broader class of goals (*Lagrangians*) for the players

Example

- i) Each P_i separately minimizes the Lagrangian

$$L_i(P) = \beta E(h_i) - S(P)$$

for some appropriate function S (e.g., entropy ...)

- ii) $\beta < \infty$ is bounded rationality

BOUNDED RATIONALITY AND COST OF COMPUTATION

- Choose $S(q) = \sum_i \int dz_i S_i(P_i(z_i))$ (e.g., entropy).

Then bounded rationality is identical to conventional, full rationality — every player wants to minimize expected cost. *Only now there is a new cost function:*

$$f_i(z, P_i) = \beta h_i(z) - S_i(P_i(z_i)) / P_i(z_i)$$

$-S_i(P_i(z_i)) / P_i(z_i)$ measures the computational cost to player i for calculating $P_i(z_i)$

COMBINING GAME THEORY AND STATISTICAL PHYSICS

- **Jaynes showed that all statistical physics ensembles arise from minimizing**

$$L_i(\mathbf{P}) = \beta E(h_i) - S(\mathbf{P}),$$

with S the Shannon entropy

- **Mean field theory arises if \mathbf{P} is a product distribution;
bounded rational game theory = mean field theory (!)**

**Much of the mathematics of statistical physics can
be applied to bounded rational game theory**

***EXAMPLE: GAMES WITH VARIABLE
NUMBERS OF PLAYERS***

- 1) The Grand Canonical Ensemble (GCE) of statistical physics models systems where the number of particles of various types varies stochastically.**
- 2) Use the underlying language, Probability Collectives, - which here is just Jaynesian inference - to translate the GCE into a game in which the number of players of various types can vary stochastically.**

***Intuition:* Players with “types” = particles with properties**

GAMES WITH VARIABLE NUMBERS OF PLAYERS - 2

Example 1 (microeconomics):

- i) A set of bounded rational companies,
- ii) with cost functions given by market valuations,
- iii) each of which must decide how many employees of various types to have.

Example 2 (evolutionary game theory):

- i) A set of species,
- ii) with cost functions given by fractions of total resources they consume,
- iii) each of which must “decide” how many phenotypes of various types to express.

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DISTRIBUTED ADAPTIVE CONTROL

- 1) Control of routers in a network.**
- 2) Control of robots working together to construct a spacestation.**
- 3) Control of flaplets on an aircraft wing.**
- 4) Control of signals to human teams performing a joint task.**
- 5) Control of variables in a parallel computer algorithm to optimize a function.**

Must be adaptive (i.e., not wed to a system model) to

- i) Avoid brittleness;***
- ii) Scale well;***
- iii) Be fault-tolerant;***
- iv) Be widely applicable, with minimal (or even no) hand-tuning.***

WHAT IS DISTRIBUTED CONTROL?

1) *A set of N agents: Joint move $x = (x_1, x_2, \dots, x_N)$*

2) *Since they are distributed, their joint probability is a product distribution:*

$$q(\mathbf{x}) = \prod_i q_i(x_i)$$

- *This definition of distributed agents is adopted from (extensive form) noncooperative game theory.*

3) *Distributed control is a common payoff game - a bounded rational (statistical physics) one.*

EXAMPLE: KSAT

- $x = \{0, 1\}^N$
- **A set of many disjunctions, “clauses”, each involving K bits.**
E.g., $(x_2 \vee x_6 \vee \sim x_7)$ is a clause for $K = 3$
- **Goal: Find a bit-string x that simultaneously satisfies all clauses. $G(x)$ is #violated clauses.**
- **For us, this goal becomes: find mixed strategy $q(x) = \prod_i q_i(x_i)$ tightly centered about such an x .**

The canonical computationally difficult problem

ITERATIVE DISTRIBUTED CONTROL

- 1) s is current uncertainty of what x to pick, i.e., uncertainty of where $q(x)$ is concentrated.
 - Early in the control process, high uncertainty.
- 2) Find q minimizing $E_q(G)$ while consistent with s .
- 3) Reduce s . Return to (2).
- 4) Stop at mixed strategy q with good (low) $E_q(G)$.

Can do (2) \rightarrow (3) without ever explicitly specifying s

ITERATIVE DISTRIBUTED CONTROL - 2

1) The central step is to “find the q that has lowest $E_q(G)$ while consistent with $S(q) = s$ ”.

2) So we must find the critical point of the Lagrangian

$$L(q, T) = E_q(G) + T[s - S(q)] ,$$

i.e., find the q and T such that $\partial L/\partial q = \partial L/\partial T = 0$

- Deep connections with statistical physics (L is “free energy” in mean-field theory), economics

3) Then we reduce s ; repeat (find next critical point).

EXAMPLE: KSAT

$$1) S(q) = -\sum_i [b_i \ln(b_i) + (1 - b_i) \ln(1 - b_i)]$$

where b_i is $q_i(x_i = \text{TRUE})$

$$2) E_q(G) = \sum_{clauses\ j, x} q(x) K_j(x)$$

$$= \sum_{clauses\ j, x, i} \prod_i q_i(x_i) K_j(x)$$

where $K_j(x) = 1$ iff x violates clause j

Our algorithm: i) Find q minimizing $E_q(G) - TS(q)$;
ii) Lower T and return to (i).

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DISTRIBUTED SEARCH FOR q

So control reduces to finding q such that $\partial L / \partial q = 0$

- 1) Since the agents make their moves in a distributed way, that q is a product distribution.**
- 2) But they must also find that q in a distributed way.**
- 3) There are two cases to consider:**
 - i) Know functional form of G .**
 - ii) Don't know functional form of G - must sample.**

MINIMIZING $L(q)$ VIA GRADIENT DESCENT

- 1) Each i works to minimize $L(q_i, q_{(i)})$ using only partial information of the other agents' distribution, $q_{(i)}$.
- 2) The $q_i(x_i)$ component of $\nabla L(q)$, projected onto the space of allowed $q_i(x_i)$, is

$$\frac{\beta E_{q_{(i)}}(G | x_i) + \ln(q_i(x_i))}{\int dx'_i [\beta E_{q_{(i)}}(G | x_i) + \ln(q_i(x'_i))]}$$

- The subtracted term ensures q stays normalized

GRADIENT DESCENT - 2

- 3) Each agent i knows its value of $\ln(q_i(x_i))$.
- 4) Each agent i knows the $E_{q(i)}(G | x_i)$ terms.

*Each agent knows how it should change
its q_i under gradient descent over $L(q)$*

- 5) Gradient descent, even for categorical variables (!), and done in a distributed way.
- 6) Similarly the Hessian can readily be estimated (for Newton's method), etc.

EXAMPLE: KSAT

- 1) Evaluate $\mathbb{E}_{q(i)}(G \mid x_i)$ - the expected number of violated clauses if bit i is in state x_i - for every i, x_i
- 2) In gradient descent, decrease each $q_i(x_i)$ by
$$\alpha[\mathbb{E}_{q(i)}(G \mid x_i) + T \ln[q_i(x_i)] - \text{const}_j]$$
 where α is the stepsize, and const_j is an easy-to-evaluate normalization constant.
- 3) We actually have a different T for each clause, and adaptively update all of them.

ADAPTIVE DISTRIBUTED CONTROL

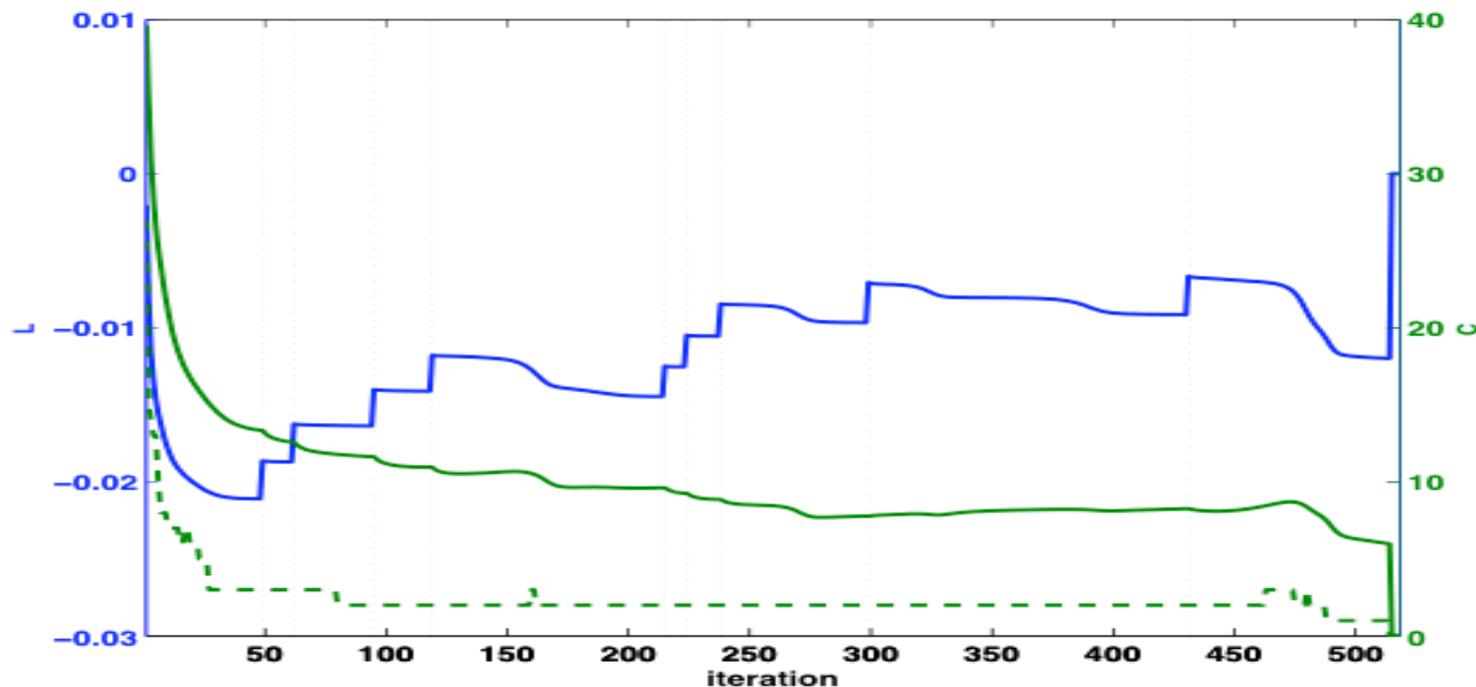
1) In *adaptive* control, don't know functional form of $G(x)$. So use Monte Carlo:

- Sample $G(x)$ repeatedly according to q ;**
- Each i independently estimates $E_{q(i)}(G | x_i)$ for all its moves x_i ;**
- Only 1 MC process, no matter how many agents**

So each q_i can adaptively estimate its update

EXAMPLE: KSAT

- i) Top plot is Lagrangian value vs. iteration;
- ii) Middle plot is average (under q) number of constraint violations;
- iii) Bottom plot is mode (under q) number of constraint violations.



CONCLUSION

- 1) Information theory - statistical inference - shows how to quantify bounded rationality*
- 2) The same mathematics underlies statistical physics; the two are identical.*
- 3) That mathematics also underlies adaptive distributed control; all three are identical.*